

# Quantum Computing and Programming Workshop Rigetti Computing 

02.25.2018


rigetti
Full-stack quantum computing company.

8-qubit and 19-qubit QPUs released on our cloud platform in 2017; 16-qubit Aspen-series QPUs released fall 2018

Quantum Cloud Services launched Fall 2018, with roadmap to 128-qubit systems

100+ employees w/ \$119M raised

Home of Fab-1, the world's first commercial quantum integrated circuit fab

Located in Berkeley, Calif. (R\&D Lab) and Fremont, Calif.

## Quantum Cloud Services and the Rigetti Forest SDK



奋

Rigetti Quantum Cloud Services

## API Model

QCS



Colocation to reduce network latency between QPU and CPU.

Parameterized program compilation to reduce \# of round trips to achieve solution.

Active reset of qubits in QPU to accelerate quantum runtime
new Quantum Machine Image access model
~30x faster than web API access models

## Web API


E.g. IBM Quantum Experience Rigetti Forest 1.x

Colocation
to reduce network latency between QPU and CPU.

+ Parametric Compilation + Active Reset
to reduce \# of round trips to achieve solution.
$\times 13.5$
$\times 34.6$ faster


## Rigetti Forest SDK

## pyQuil

intuitive Python library for writing Quil programs to run on Forest SDK.

## Forest SDK

QVM: local simulator on up to 26 qubits

QUILC: compiler with the ability to optimize programs to different architectures

## pyQuil



## Why build a

 quantum computer?


## Quantum Algorithms' Progression

1992-4
First Quantum Algorithms w/ Exponential Speedup (Deutsch-Jozsa, Shor's Factoring, Discrete Log, ...)

1996
First Quantum Database Search Algorithm (Grover's)

## 2007

Quantum Linear Equation Solving (Harrow, Hassidim, Lloyd) 2008
Quantum Algorithms for SVM's \& Principal Component Analysis

## 2013

Practical Quantum Chemistry Algorithms (VQE)
2016
Practical Quantum Optimization Algorithms (QAOA)
Simulations on Near-term Quantum Supremacy

These algorithms require big, perfect quantum computers
$>10,000,000$ qubits for Shor's algorithms
to factor a 2048 bit number

## Hybrid quantum/classical algorithms


"Quantum computing in the NISQ era and beyond" Preskill, 2018 https://arxiv.org/abs/1801.00862

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Machine Learning

Development of new training sets and algorithms

Classification and sampling of large data sets


Supply Chain
Optimization

Forecast and optimize for future inventory demand

NP-hard scheduling and logistics map into quantum applications


## Robotic <br> Manufacturing

Reduce
manufacturing time and cost

Maps to a Traveling Salesman Problem addressable by quantum constrained optimization


## Computational Materials Science

Design of better catalysts for batteries

Quantum algorithms for calculating electronic structure

Alternative Energy Research

Efficiently convert atmospheric CO2 to methanol

Powered by existing hybrid quantum- classical algorithms + machine learning

## Who is building and investing in quantum computers?

Investments across academia, government, and industry are global and growing

## No small effort



## Bits vs. Probabilistics Bits vs. Qubits

Bits Probabilistic Bits Qubits

| State (single unit) | Bit $\in\{0,1\}$ |
| :--- | :--- |


| Bits | Probabilistic Bits |  |
| :--- | :--- | :--- | :--- |
| State (single unit) | Bit $\in\{0,1\}$ | Real vector  <br> $\vec{v}=a \overrightarrow{0}+b \overrightarrow{1}$ $a, b \in \mathbb{R}_{+}$ <br>   <br>   |



| Bits |  | Probabilistic Bits |  | Qubits |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| State (single unit) | Bit $\in\{0,1\}$ | Real vector $\vec{v}=a \overrightarrow{0}+b \overrightarrow{1}$ | $\begin{aligned} & a, b \in \mathbb{R}_{+} \\ & a+b=1 \end{aligned}$ | Complex vector $\|\psi\rangle=\alpha\|0\rangle+\beta\|1\rangle$ | $\begin{gathered} \alpha, \beta \in \mathbb{C} \\ \|\alpha\|^{2}+\|\beta\|^{2}=1 \end{gathered}$ |


| Bits |  | Probabilistic Bits |  | Qubits |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| State (single unit) | Bit $\in\{0,1\}$ | Real vector $\vec{v}=a \overrightarrow{0}+b \overrightarrow{1}$ | $\begin{aligned} & a, b \in \mathbb{R}_{+} \\ & a+b=1 \end{aligned}$ | Complex vector $\|\psi\rangle=\alpha\|0\rangle+\beta\|1\rangle$ | $\begin{gathered} \alpha, \beta \in \mathbb{C} \\ \|\alpha\|^{2}+\|\beta\|^{2}=1 \end{gathered}$ |
|  |  |  | $\|\alpha\|^{2}=\operatorname{Pr}$ |  | Probability of 1 |

## Bits vs. Probabilistic Bits vs. Qubits

| Bits | Probabilistic Bits |  | Qubits |  |  |
| :--- | :--- | :--- | ---: | ---: | :---: |
| State (single unit) | Bit $\in\{0,1\}$ | Real vector | $a, b \in \mathbb{R}_{+}$ | Complex vector | $\alpha, \beta \in \mathbb{C}$ |
|  |  | $\vec{v}=a \overrightarrow{0}+b \overrightarrow{1}$ | $a+b=1$ | $\|\psi\rangle=\alpha\|0\rangle+\beta\|1\rangle$ | $\|\alpha\|^{2}+\|\beta\|^{2}=1$ |

$$
\vec{v}_{\text {coin }}=\frac{1}{2} \overrightarrow{0}+\frac{1}{2} \overrightarrow{1}
$$

## Bits vs. Probabilistic Bits vs. Qubits

Bits
Probabilistic Bits
Qubits

| State (single unit) | Bit $\in\{0,1\}$ | Real vector | $a, b \in \mathbb{R}_{+}$ | Complex vector | $\alpha, \beta \in \mathbb{C}$ |
| :--- | :--- | ---: | ---: | ---: | ---: |
|  | $\vec{v}=a \overrightarrow{0}+b \overrightarrow{1}$ | $a+b=1$ | $\|\psi\rangle=\alpha\|0\rangle+\beta\|1\rangle$ | $\|\alpha\|^{2}+\|\beta\|^{2}=1$ |  |

CLASSICAL BIT

$$
\vec{v}_{\mathrm{coin}}=\frac{1}{2} \overrightarrow{0}+\frac{1}{2} \overrightarrow{1}
$$

$$
\begin{aligned}
& |\psi\rangle_{\mathrm{coin}}=\frac{1}{\sqrt{2}}|0\rangle+\frac{1}{\sqrt{2}}|1\rangle \\
& |\psi\rangle_{\mathrm{coin}}=\frac{1}{\sqrt{2}}|0\rangle-\frac{1}{\sqrt{2}}|1\rangle \\
& |\psi\rangle_{\mathrm{coin}}=\frac{1}{\sqrt{2}}|0\rangle-\frac{i}{\sqrt{2}}|1\rangle
\end{aligned}
$$

## Bits vs. Probabilistic Bits vs. Qubits

Bits
Probabilistic Bits
Qubits

| State (single unit) | Bit $\in\{0,1\}$ | Real vector | $a, b \in \mathbb{R}_{+}$ | Complex vector | $\alpha, \beta \in \mathbb{C}$ |
| :--- | :--- | ---: | ---: | ---: | ---: |
|  | $\vec{v}=a \overrightarrow{0}+b \overrightarrow{1}$ | $a+b=1$ | $\|\psi\rangle=\alpha\|0\rangle+\beta\|1\rangle$ | $\|\alpha\|^{2}+\|\beta\|^{2}=1$ |  |

$$
\vec{v}_{\text {coin }}=\frac{1}{2} \overrightarrow{0}+\frac{1}{2} \overrightarrow{1} \quad|\psi\rangle_{\text {coin }}=\frac{1}{\sqrt{2}}|0\rangle+\frac{e^{i \theta}}{\sqrt{2}}|1\rangle
$$

CLASSICAL BIT

## Bits vs. Probabilistic Bits vs. Qubits

Bits
Probabilistic Bits
Qubits

| State (single unit) | Bit $\in\{0,1\}$ | Real vector | $a, b \in \mathbb{R}_{+}$ | Complex vector | $\alpha, \beta \in \mathbb{C}$ |
| :--- | :--- | ---: | ---: | ---: | ---: |
|  | $\vec{v}=a \overrightarrow{0}+b \overrightarrow{1}$ | $a+b=1$ | $\|\psi\rangle=\alpha\|0\rangle+\beta\|1\rangle$ | $\|\alpha\|^{2}+\|\beta\|^{2}=1$ |  |



## Qubit States



$$
|\psi\rangle=\cos \left(\frac{\theta}{2}\right)|0\rangle+e^{i \varphi} \sin \left(\frac{\theta}{2}\right)|1\rangle
$$

Qubit:

$$
|\psi\rangle=\alpha|0\rangle+\beta|1\rangle, \quad \alpha, \beta \in \mathbb{C}, \quad|\alpha|^{2}+|\beta|^{2}=1
$$

Kets:

$$
|0\rangle=\binom{1}{0} \quad|1\rangle=\binom{0}{1} \quad|\psi\rangle=\binom{\alpha}{\beta}
$$

## Measurement yields:

- '0' with probability $|\alpha|^{2}$
- ' 1 ' with probability $|\beta|^{2}$

Qubit:

$$
|\psi\rangle=\alpha|0\rangle+\beta|1\rangle, \quad \alpha, \beta \in \mathbb{C}, \quad|\alpha|^{2}+|\beta|^{2}=1
$$

Kets:

Bras:

$$
|0\rangle=\binom{1}{0} \quad|1\rangle=\binom{0}{1} \quad|\psi\rangle=\binom{\alpha}{\beta}
$$

$$
\langle 0|=(1,0)
$$

$$
\langle 1|=(0,1)
$$

$$
\langle\psi|=(\bar{\alpha}, \bar{\beta})
$$

## Brackets

(Inner Product):

$$
\begin{gathered}
|\phi\rangle=\gamma|0\rangle+\delta|1\rangle \\
\langle\phi \mid \psi\rangle=\bar{\gamma} \alpha+\bar{\delta} \beta=\overline{\langle\psi \mid \phi\rangle}
\end{gathered}
$$

Multiple qubits:

$$
|\psi\rangle_{n-1} \otimes \ldots \otimes|\psi\rangle_{2} \otimes|\psi\rangle_{1} \otimes|\psi\rangle_{0}
$$

Tensor product:

$$
\begin{aligned}
&|\psi\rangle \otimes|\phi\rangle=\left(\alpha_{0}|0\rangle+\alpha_{1}|1\rangle\right) \otimes\left(\beta_{0}|0\rangle+\beta_{1}|1\rangle\right) \\
&=\alpha_{0} \beta_{0}|00\rangle+\alpha_{0} \beta_{1}|01\rangle+\alpha_{1} \beta_{0}|10\rangle+\alpha_{1} \beta_{1}|11\rangle
\end{aligned}
$$

Vector form: $\quad\binom{\alpha_{0}}{\alpha_{1}} \otimes\binom{\beta_{0}}{\beta_{1}}=\left(\begin{array}{c}\alpha_{0}\left(\begin{array}{c}\beta_{0} \\ \beta_{1} \\ \beta_{0} \\ \beta_{0} \\ \beta_{1}\end{array}\right)\end{array}\right)=\left(\begin{array}{c}\alpha_{0} \beta_{0} \\ \alpha_{0} \beta_{1} \\ \alpha_{1} \beta_{0} \\ \alpha_{1} \beta_{1}\end{array}\right)$

## Associative:

$$
(|\psi\rangle \otimes|\phi\rangle) \otimes|\gamma\rangle=|\psi\rangle \otimes(|\phi\rangle \otimes|\gamma\rangle)
$$

Not commutative:

$$
|\psi\rangle \otimes|\phi\rangle \neq|\phi\rangle \otimes|\psi\rangle
$$

# Single Qubit Operations 

Unitary Operators (Gates):

$$
U U^{\dagger}=U^{\dagger} U=I
$$

$$
|\psi\rangle \rightarrow\left|\psi^{\prime}\right\rangle=U|\psi\rangle
$$

Preserve inner product:

$$
\begin{gathered}
\langle\phi| \rightarrow\left\langle\phi^{\prime}\right|=\langle\phi| U^{\dagger} \\
\langle\phi \mid \psi\rangle \rightarrow\left\langle\phi^{\prime} \mid \psi^{\prime}\right\rangle=\langle\phi| U^{\dagger} U|\psi\rangle=\langle\phi \mid \psi\rangle
\end{gathered}
$$

$$
\begin{array}{ll}
I=\left(\begin{array}{ll}
1 & 0 \\
0 & 1
\end{array}\right) & Y=\left(\begin{array}{cc}
0 & -i \\
i & 0
\end{array}\right) \\
X=\left(\begin{array}{ll}
0 & 1 \\
1 & 0
\end{array}\right) & Z=\left(\begin{array}{cc}
1 & 0 \\
0 & -1
\end{array}\right)
\end{array}
$$

$$
\begin{array}{ll}
I|0\rangle=|0\rangle & \left(\begin{array}{ll}
1 & 0 \\
0 & 1
\end{array}\right)\binom{0}{1}=\binom{0}{1} \\
I|1\rangle=|1\rangle & \left(\begin{array}{ll}
1 & 0 \\
0 & 1
\end{array}\right)\binom{1}{0}=\binom{1}{0}
\end{array}
$$

Qubit Operations: Pauli-X (NOT) gate

$$
\begin{array}{ll}
X|0\rangle=|1\rangle & \left(\begin{array}{ll}
0 & 1 \\
1 & 0
\end{array}\right)\binom{1}{0}=\binom{0}{1} \\
X|1\rangle=|0\rangle & \left(\begin{array}{ll}
0 & 1 \\
1 & 0
\end{array}\right)\binom{0}{1}=\binom{1}{0}
\end{array}
$$

$$
\begin{array}{ll}
Y|0\rangle=i|1\rangle & \left(\begin{array}{cc}
0 & -i \\
i & 0
\end{array}\right)\binom{1}{0}=i\binom{0}{1} \\
Y|1\rangle=-i|0\rangle & \left(\begin{array}{cc}
0 & -i \\
i & 0
\end{array}\right)\binom{0}{1}=-i\binom{1}{0}
\end{array}
$$

$$
\begin{array}{ll}
Z|0\rangle=|0\rangle & \left(\begin{array}{cc}
1 & 0 \\
0 & -1
\end{array}\right)\binom{1}{0}=\binom{1}{0} \\
Z|1\rangle=-|1\rangle & \left(\begin{array}{cc}
1 & 0 \\
0 & -1
\end{array}\right)\binom{0}{1}=-\binom{0}{1}
\end{array}
$$

$$
H=\frac{1}{\sqrt{2}}\left(\begin{array}{cc}
1 & 1 \\
1 & -1
\end{array}\right)
$$

$$
\begin{aligned}
& H|0\rangle=\frac{1}{\sqrt{2}}(|0\rangle+|1\rangle)
\end{aligned} \frac{1}{\sqrt{2}}\left(\begin{array}{cc}
1 & 1 \\
1 & -1
\end{array}\right)\binom{1}{0}=\frac{1}{\sqrt{2}}\binom{1}{1}
$$

## Multi-Qubit Operations

## Examples:

$$
I \otimes X(|0\rangle \otimes|0\rangle)=|0\rangle \otimes|1\rangle=I_{1} X_{0}|00\rangle=|01\rangle
$$

$$
X \otimes H(|0\rangle \otimes|0\rangle)=|1\rangle \otimes \frac{1}{\sqrt{2}}(|0\rangle+|1\rangle)=X_{1} H_{0}|00\rangle=\frac{1}{2}(|10\rangle+|11\rangle)
$$

$$
\begin{aligned}
& A \otimes B=\left(\begin{array}{ll}
A_{00} & A_{01} \\
A_{10} & A_{11}
\end{array}\right) \otimes\left(\begin{array}{ll}
B_{00} & B_{01} \\
B_{10} & B_{11}
\end{array}\right)=\left(\begin{array}{lll}
A_{00}[B] & A_{01}[B] \\
A_{10}[B] & A_{11}[B]
\end{array}\right) \\
&=\left(\begin{array}{llll}
A_{00} B_{00} & A_{00} B_{01} & A_{01} B_{00} & A_{01} B_{01} \\
A_{00} B_{10} & A_{00} B_{11} & A_{01} B_{10} & A_{01} B_{11} \\
A_{10} B_{00} & A_{01} B_{01} & A_{11} B_{00} & A_{11} B_{01} \\
A_{10} B_{10} & A_{11} B_{11} & A_{11} B_{10} & A_{11} B_{11}
\end{array}\right)
\end{aligned}
$$

$X|0\rangle=|1\rangle$

$X_{1} H_{0}|00\rangle$


$$
c U=|0\rangle\langle 0| \otimes I+|1\rangle\langle 1| \otimes U
$$

If qubit 1 is in the state $\mid 0>$, apply $I$ (identity) to qubit 0 Else if qubit 1 is in the state $\mid 1>$, apply $U$ to qubit 0

For example,

$$
\text { CNOT }=|0\rangle\langle 0| \otimes I+|1\rangle\langle 1| \otimes X=\left(\begin{array}{llll}
1 & 0 & 0 & 0 \\
0 & 1 & 0 & 0 \\
0 & 0 & 0 & 1 \\
0 & 0 & 1 & 0
\end{array}\right)
$$

A state that cannot be written as a product state, i.e.

$$
|\psi\rangle \neq|\xi\rangle \otimes|\phi\rangle
$$

An example of a state that is not entangled:

$$
\frac{1}{\sqrt{2}}(|00\rangle+|01\rangle)=|0\rangle \otimes \frac{1}{\sqrt{2}}(|0\rangle+|1\rangle)
$$

An example of a state that is entangled:

$$
\frac{1}{\sqrt{2}}(|00\rangle+|11\rangle)
$$

Project a ket/bra along a given ket/bra via its corresponding projection operator.

For some $|\psi\rangle$ the corresponding projection operator is given by the outer product

$$
P_{\psi}=|\psi\rangle\langle\psi|=\binom{\alpha}{\beta}(\bar{\alpha}, \bar{\beta})=\left(\begin{array}{cc}
|\alpha|^{2} & \alpha \bar{\beta} \\
\beta \bar{\alpha} & |\beta|^{2}
\end{array}\right)
$$

e.g.

$$
P_{0}=\left(\begin{array}{ll}
1 & 0 \\
0 & 0
\end{array}\right), P_{1}=\left(\begin{array}{ll}
0 & 0 \\
0 & 1
\end{array}\right), P_{+}=\frac{1}{2}\left(\begin{array}{ll}
1 & 1 \\
1 & 1
\end{array}\right)
$$

Computational basis: $\quad\{|0\rangle,|1\rangle\}$

Measurement yields:

- '0' with probability

$$
\langle\psi| P_{0}|\psi\rangle=\langle\psi \mid 0\rangle\langle 0 \mid \psi\rangle=|\langle 0 \mid \psi\rangle|^{2}
$$

- ' 1 ' with probability

$$
\langle\psi| P_{1}|\psi\rangle=\langle\psi \mid 1\rangle\langle 1 \mid \psi\rangle=|\langle 1 \mid \psi\rangle|^{2}
$$

Some other basis: $\quad\left\{\left|0^{\prime}\right\rangle=U|0\rangle,\left|1^{\prime}\right\rangle=U|1\rangle\right\}$

$$
\left|\psi^{\prime}\right\rangle=U^{\dagger}|\psi\rangle
$$

Measurement yields:

- $0^{\prime}$ with probability

$$
\begin{gathered}
\langle\psi| P_{0^{\prime}}|\psi\rangle=\left\langle\psi \mid 0^{\prime}\right\rangle\left\langle 0^{\prime} \mid \psi\right\rangle=\langle\psi| U|0\rangle\langle 0| U^{\dagger}|\psi\rangle \\
=\left\langle\psi^{\prime} \mid 0\right\rangle\left\langle 0 \mid \psi^{\prime}\right\rangle=\left|\left\langle 0 \mid \psi^{\prime}\right\rangle\right|^{2}
\end{gathered}
$$

- 1 ' with probability

$$
\begin{gathered}
\langle\psi| P_{1^{\prime}}|\psi\rangle=\left\langle\psi \mid 1^{\prime}\right\rangle\left\langle 1^{\prime} \mid \psi\right\rangle=\langle\psi| U|1\rangle\langle 1| U^{\dagger}|\psi\rangle \\
=\left\langle\psi^{\prime} \mid 1\right\rangle\left\langle 1 \mid \psi^{\prime}\right\rangle=\left|\left\langle 1 \mid \psi^{\prime}\right\rangle\right|^{2}
\end{gathered}
$$

Measurement of $\mid \Psi>$ in some basis $\{\mathrm{U}|0>, \mathrm{U}| 1>\}$
$=$ Measurement of $\mathrm{U}^{\dagger} \mid \Psi>$ in standard computational basis $\{|0>| 1>$,

## Quantum Programs

$X|0\rangle=|1\rangle$

from pyquil import Program, get_qc
from pyquil.gates import X
p = Program (X(0))
qc = get_qc('9q-generic-qvm')
results = qc.run_and_measure $(p$, trials=10)[0]
print (results)
[1 1111111111 ]
$X|0\rangle=|1\rangle$
from pyquil import Program, get_qc
from pyquil.gates import X, MEASURE
p = Program()
p.declare('ro', 'BIT', 1)
p.inst(X(0))
p.inst(MEASURE(0, 'ro'))
p.wrap_in_numshots_loop(shots=10)
qc = get_qc('9q-generic-qvm')
results = qc.run(qc.compile(p))
print (p)


DECLARE ro BIT[1]
X 0
MEASURE 0 ro[0]

$$
X|0\rangle=|1\rangle
$$

from pyquil import Program, get_qc
from pyquil.gates import X, MEASURE
p = Program(
p.declare('ro', 'BIT', 1)
p.inst(X(0))
p.inst(MEASURE(0, 'ro'))
p.wrap_in_numshots_loop(shots=10)
qc = get_qc('9q-generic-qvm')
results = qc.run(qc.compile(p))
print (results)
[1][1][1][1][1][1]

## $X_{1} H_{0}|0\rangle_{1}|0\rangle_{0}$

from pyquil import Program, get_qc from pyquil.gates import X, H, MEASURE
p = Program(
ro = p.declare('ro', 'BIT', 2)
p.inst( $\mathrm{H}(0)$ )
p.inst(X(1))
p.inst(MEASURE(0, ro[0]))
p.inst(MEASURE(1, ro[1]))
p.wrap_in_numshots_loop(shots=10)
qc = get_qc('9q-generic-qvm')
results = qc.run(qc.compile(p))
print (results)

[ $\left[\begin{array}{ll}0 & 1\end{array}\right]$
[1 1]
[01]
[11]
[11]
[01]
[11]
[01]
$\left[\begin{array}{ll}0 & 1\end{array}\right]$
[0 1]]

## $Y_{1} Z_{0} X_{1} H_{0}|0\rangle_{1}|0\rangle_{0}$

from pyquil import Program, get_qc
from pyquil.gates import X, Y, Z, H, MEASURE

[[1 0]
$\mathrm{p}=$ Program() [1 0]
ro = p.declare('ro', 'BIT', 2)
$\mathrm{p}+=\operatorname{Program}(\mathrm{H}(0), \mathrm{X}(1), \mathrm{Z}(0), \mathrm{Y}(1), \operatorname{MEASURE}(0, \mathrm{ro}[0])$, MEASURE(1, ro[1])) [0 0] p.wrap_in_numshots_loop(shots=10)
$\mathrm{qc}=$ get_qc('9q-generic-qvm')
results = qc.run(qc.compile(p))
[10] [1 0]
[0 0]
print (results)
[10]
[1 0]
[0 0]
[0 0 0]]


## Quantum Programming Examples

## Quantum Die Example

## Goal: Create a fair N -sided die

## Question: What gate would we use?

$$
H^{\otimes n}|0\rangle^{\otimes n}=(H|0\rangle)^{\otimes n}=\frac{1}{\sqrt{2^{n}}} \sum_{z=0}^{2^{n}-1}|z\rangle
$$

# Question: How many qubits would we use? 

$$
2^{n}=N \Rightarrow n=\log _{2} N
$$

Quantum
Teleportation
Example

Goal: Teleport a qubit state from Alice to Bob

Goal: Teleport a qubit state from Alice to Bob

Scenario: Alice is in possession of a qubit | $\Psi\rangle$, which she would like to teleport over to Bob, who is at some distant location.

## Protocol:

- Create a Bell state, giving one qubit each to Alice and Bob
- Have Alice measure both her qubits in the Bell basis, and send her results to Bob
- Have Bob conditionally apply gates to his qubits, based off Alice's measurements, to reconstruct the original qubit at his location

from pyquil import Program
from pyquil.gates import $I, X$
from pyquil.api import WavefunctionSimulator
$p=\operatorname{Program}(X(0))$
ro = p.declare('ro', 'BIT', 1)
p.measure(0, ro[0]).if_then(ro[0], Program(X(1)), Program(I(1)))
wfn = WavefunctionSimulator().wavefunction(p)
print (wfn)
(1+0j)|11>
from pyquil import Program
from pyquil.gates import $I, X$
from pyquil.api import WavefunctionSimulator
p = Program $(1(0))$
ro = p.declare('ro', 'BIT', 1)
p.measure(0, ro[0]).if_then(ro[0], Program(X(1)), Program(I(1)))
wfn = WavefunctionSimulator().wavefunction(p)
print (wfn)
$(1+0 \mathrm{j}) \mid 00>$


## Example: Quantum Teleportation



## Quantum

 Approximate Optimization Algorithm (QAOA)Goal: Given binary constraints over bitstrings

$$
\begin{aligned}
& z \in\{0,1\}^{n} \\
& C_{\alpha}(z)= \begin{cases}1 & \text { if } z \text { satisfies the constraint } \alpha \\
0 & \text { if } z \text { does not }\end{cases}
\end{aligned}
$$

Find the bitstring that maximizes the objective function

$$
\operatorname{argmax}_{z} C(z)=\operatorname{argmax}_{z} \sum_{\alpha=1}^{m} C_{\alpha}(z)
$$

## MaxCut problem:

Given some undirected graph with arbitrary (non-negative) weights, find a partition $(S, \bar{S})$ of the graph's nodes (a 'cut' of the graph) that maximizes the weights along the cut

$$
\sum_{i \in S, j \in \bar{S}} w_{i j}
$$




MaxCut solution (as a bitstring):

$$
01001 \text { OR } 10110
$$

On a quantum computer:

$$
\frac{1}{\sqrt{2}}(|01001\rangle+|10110\rangle)
$$

(ideally)


MaxCut objective function:

$$
\sum_{i \in S, j \in \bar{S}} w_{i j}
$$

On a quantum computer:

$$
\frac{1}{2} \sum_{i, j \in V} w_{i j} \frac{1-Z_{i} Z_{j}}{2}
$$

## Noise and Quantum Computation

'Pure' quantum states

$$
|\psi\rangle=\alpha|0\rangle+\beta|1\rangle
$$

evolve via Unitary operations

$$
|\psi\rangle \rightarrow\left|\psi^{\prime}\right\rangle=U|\psi\rangle \quad U U^{\dagger}=U^{\dagger} U=I
$$

More generally, quantum states are described by "Density Matrix"

$$
\rho=\sum_{i} p_{i}\left|\psi_{i}\right\rangle\left\langle\psi_{i}\right|
$$

evolving via Kraus operations ("quantum channel")

$$
\rho \rightarrow \sum_{i} K_{i} \rho K_{i}^{\dagger}, \quad \sum_{i} K_{i}^{\dagger} K_{i}=I
$$

For example,

$$
\rho=\frac{1}{2}(|0\rangle\langle 0|+|1\rangle\langle 1|)=\frac{1}{2}\left(\begin{array}{ll}
1 & 0 \\
0 & 1
\end{array}\right)
$$

Not to be confused with

$$
\rho=\left(\frac{1}{\sqrt{2}}(|0\rangle+|1\rangle)\right)\left(\frac{1}{\sqrt{2}}(\langle 0|+\langle 1|)\right)=\frac{1}{2}\left(\begin{array}{ll}
1 & 1 \\
1 & 1
\end{array}\right)
$$

Example of quantum channel/set of Kraus operators/noise model:

$$
\{\sqrt{p} X, \sqrt{1-p} Z\}
$$

Quantum state passing through the channel/experiencing the noise transforms to:

$$
\rho \rightarrow p X^{\dagger} \rho X+(1-p) Z^{\dagger} \rho Z
$$



## Thank you, and keep in touch! <br> amy@rigetti.com

## Join our Slack channel: rigetti-forest.slack.com

